ON A FUNCTIONAL EQUATION IN TWO AND THREE VARIABLES WITH MEASURES OF KIND β IN INFORMATION THEORY

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(Received 4 October 1973)

Starting from a functional equation of two variables containing a parameter β , the information function of kind β given have been obtained. The paper contains a characterization and properties of this measure. Also the study has been extended to the information improvement of the prediction probability distribution (q_1, q_2, \ldots, q_n) revised as $(r_1, .r_2, \ldots, r_n)$ on the basis of realization (p_1, p_2, \ldots, p_n) .

Let
$$P = (p_1, p_2, ..., p_n)$$
 $p_i \ge 0, \sum_{i=1}^n p_i = 1$ be a probability distribution of a set of *n* events

on the basis of an experiment E whose true (prediction) probability distribution is $Q = (q_1, q_2, \ldots, q_n)$,

 $q_i \ge 0$, $\sum_{i=1}^n q_i = 1$, then Kullback's¹ measure of information that P provides about Q is

$$I(P;Q) = \sum_{i=1}^{n} p_i \log (p_i q_i^{-1}).$$
 (1)

Again let $R = (r_1, r_2, ..., r_n), r_i \ge 0, \sum_{i=1}^n r_i = 1$ be a revised probability distribution of a

original probability distribution $Q = (q_1, q_2, ..., q_n), q_i \ge 0, \sum_{i=1}^n q_i = 1$ on the basis of realization

 $P = (p_1, p_2, ..., p_n), p_i \ge 0, \sum_{i=1}^n p_i = 1$ obtained from the experiment *E*. Then the information improvement² is given by

$$I(P;Q;R) = \sum_{i=1}^{n} p_i \log (r_i q_i^{-1}).$$
(2)

In this paper we shall take the functional equation in two and three variables, involving a parameter β which under suitable boundary conditions would give the new measures named as the 'information' and 'information-improvement' of kind β^3 . The paper contains the characterizations of these new measures and some of their properties.

Sharma and Autar⁴ have studied earlier^{5,6} some other generalizations of (1) and (2) differently.

CHARACTERIZATION OF INFORMATION OF KIND &

We shall first take generalization of (1).

Let the information of kind β of P with respect to Q denoted by L_n be taken to satisfy the following postulates :

Postulate I

$$L_{n}\left(\begin{array}{c}p_{1}, p_{2}, \ldots, p_{n}\\q_{1}, q_{2}, \ldots, q_{n}\end{array}\right) = L_{n-1}\left(\begin{array}{c}P_{2}, p_{3}, \ldots, p_{n}\\Q_{2}, q_{3}, \ldots, q_{n}\end{array}\right) + P_{2}^{\beta} Q_{2}^{1-\beta} L_{2}\left(\begin{array}{c}p_{1}/P_{2}, p_{2}/P_{2}\\q_{1}/Q_{2}, q_{2}/Q_{2}\end{array}\right).$$

where $P_2 = p_1 + p_2$, $Q_2 = q_1 + q_2$, P_2 , $Q_2 > 0$.

Postulate II

 $L_3\left(\frac{p_1, p_2, p_3}{q_1, q_2, q_3}\right)$ is a symmetric function such that for any permutation of p's there is the same

permutation of q's.

Postulate III

$$L_2 \left(\begin{array}{c} \beta^{1/\beta}, 1-\beta^{1/\beta} \\ 1 & , 0 \end{array} \right) = 1. \ (0 < \beta < 1).$$

Note 1: The postulate I is a recursivity of information of kind β and it is a relation corresponding to additivity. Also postulate III can be considered as a normalization property and here it includes the parameter β for the quantities were $0 < \beta < 1$.

Note 2 : If we take

$$f(x, y) = L_2 \left(\begin{array}{c} x, 1-x \\ y, 1-y \end{array} \right), \tag{3}$$

then with the help of postulates I and II, L_n $(n \ge 3)$ can be expressed in terms of the single function f(x, y).

The postulate I at once gives the elegant form

$$L_{n} = \sum_{i=2}^{n} P_{i}\beta Q_{i}^{1-\beta} \cdot f (p_{i}/P_{i}, q_{i}/Q_{i}), \qquad (4)$$

where $P_i = p_1 + p_2 + \ldots + p_i$, $Q_i = q_1 + q_2 + \ldots + q_i$, $i = 1, 2, \ldots, n$ with $P_n = Q_n = 1$.

When recursivity and symmetry of information of kind β are applied to L_3 , we obtain the functional equation

$$f(x, y) + (1-x)^{\beta} (1-y)^{1-\beta} f[u/(1-x), v/(1-y)] = f(u, v) + (1-u)^{\beta} (1-v)^{1-\beta} f[x/(1-u), y/(1-v)],$$
for x, y, u, v \epsilon [0, 1] with x + u, y + v \epsilon [0, 1].
(5)

If we put y = x and v = u and take f(x, x) = f(x), (5) reduces to Kendall's⁸ functional equation of information function.

We now adopt the following definition.

Definition: A real-valued solution f(x, y) of (5) defined on [0, 1] \times [0, 1] is said to be an information function of kind $\beta \in (0, 1)$, if it satisfies the following boundary conditions

$$f(0, 0) = f(1, 1), \tag{6}$$

$$f(\beta^{1/\beta}, 1) = f(1 - \beta^{1/\beta}, 0) = 1,$$
(7)

Again if f(x, y) is an information function of kind β then the information of kind β that P provides about Q is given by (4).

It is shown below that f(x, y) = f(1-x, 1-y) and f(0, 0) = 0. Thus (7) follows by putting $x = \beta^{1/\beta}$ and y = 1 in (5).

We now give characterizations of information function and information of kind $\beta \neq 1$ in theorems I and II respectively.

SONI : Functional Equation in Information Theory

Theorem I

The only solution
$$f(x, y)$$
 of (5) satisfying the additional conditions (6), and (7) is given by

$$f(x, y) = \begin{bmatrix} 1 - x\beta \ y^1 - \beta - (1 - x)\beta \ (1 - y)^1 - \beta \end{bmatrix} (1 - \beta)^{-1}, \quad (8)$$
for $(x, y) \in [0, 1] \times [0, 1] \sim (0, 1) U(1, 0)$; $\beta \in (0, 1)$,
(We shall take $0^{\alpha} = 0 \quad (\alpha \neq 0)$.

Proof :

Taking
$$x = y = 0$$
 in (5), $f(0, 0) = 0$ and hence from (6), we get

$$f(1, 1) = f(0, 0) = 0.$$
(9)

Now replacing u, v in (5) by 1 - x and 1 - y respectively and applying (9), we have

$$f(x, y) = f(1-x, 1-y) \text{ for } x, y \in [0, 1].$$
(10)

Let p_1 , p_2 , q_1 , q_2 be four arbitrary numbers belonging to (0, 1). Setting $p_1 = 1-x$, $p_2 = u(1-x)^{-1}$, $q_1 = 1-y$, $q_2 = v(1-y)^{-1}$ in (5) and using (10), we have

$$f(p_1, q_1) + p_1^{\beta} q_1^{1-\beta} f(p_2, q_2) = f(p_1 p_2, q_1 q_2) + (1 - p_1 p_2)^{\beta} (1 - q_1 q_2)^{1-\beta} .$$

$$f(1 - p_1 / (1 - p_1 p_2), 1 - q_1 / (1 - q_1 q_2))$$

$$(11)$$

for
$$p_1, q_1 \in [0, 1], p_2, q_2 \in [0, 1]$$
 such that $p_1 p_2 \neq 1$ and $q_1 q_2 \neq 1$.

Take the function

$$F(p_1, p_2; q_1, q_2) = f(p_1, q_1) + [p_1^{\beta} q_1^{1} - \beta + (1 - p_1)^{\beta} (1 - q_1)^{1} - \beta] f(p_2, q_2),$$

where $p_1, p_2, q_1, q_2 \in (0, 1).$ (12)

We shall show that $F(p_1, p_2; q_1, q_2)$ is symmetric, *i.e.*,

$$F(p_1, p_2; q_1, q_2) = F(p_2, p_1; q_2, q_1).$$
(13)

From (11) and (12) we have

$$F(p_1, p_2; q_1, q_2) = f(p_1 p_2, q_1 q_2) + (1 - p_1 p_2)\beta (1 - q_1 q_2)^1 - \beta [f(1 - p_1/(1 - p_1 p_2), 1 - q_1/(1 - q_1 q_2)) + (1 - p_1/(1 - p_1 p_2))\beta (1 - q_1/(1 - q_1 q_2))^1 - \beta f(p_2, q_2)]$$
(14)

Setting $p_1^* = (1 - p_1) (1 - p_1 p_2)^{-1}$, $q_1^* = (1 - q_1) (1 - q_1 q_2)^{-1}$, we have $A (p_1, p_2; q_1, q_2) = f (1 - p_1 / (1 - p_1 p_2), 1 - q_1 / (1 - q_1 q_2)) + (1 - p_1 / (1 - p_1 p_2))^{\beta} \cdot (1 - q_1 / (1 - q_1 q_2))^{1 - \beta} f (p_2, q_2)$ $= f (p_1^*, q_1^*) + p_1^{*\beta} q_1^{*(1-\beta)} f(p_2, q_2)$

using (10) and (11) we get

$$= f\left(\frac{1-p_2}{1-p_1p_2}, \frac{1-q_2}{1-q_1q_2}\right) + \left(\frac{1-p_2}{1-p_1p_2}\right)^{\beta} \left(\frac{1-q_2}{1-q_1q_2}\right)^{1-\beta} f(p_1, q_1)$$

= $A(p_2, p_1; q_2, q_1).$

So it follows from (12) that

 $F(p_1, p_2; q_1, q_2) - F(p_2, p_1; q_2, q_1) = 0$, proving (13).

Next putting $p_2 = \beta^{1/\beta}$, $q_2 = 1$ or $p_2 = 1 - \beta^{1/\beta}$, $q_2 = 0$ in (13) and using the definition of $F(p_1, p_2; q_1, q_2)$ and (7) we get

$$0 = f(p_1, q_1) + [p_1\beta q_1^1 - \beta + (1 - p_1)\beta (1 - q_1)^1 - \beta] - 1 - \beta f(p_1, q_1)$$

From this it follows that

$$f(p_1, q_1) = [1 - p_1^{\beta} q_1^{1-\beta} - (1 - p_1)^{\beta} (1 - q_1)^{1-\beta}] (1 - \beta)^{-1}, \text{ for all } p_1, q_1 \in (0, 1)$$

Since f(0, 0) = f(1, 1) = 0, the result is true when $p_1 = q_1 = 0$ or $p_1 = q_1 = 1$.

Theorem II

Let $P = (p_1, p_2, \ldots, p_n)$, $Q = (q_1, q_2, \ldots, q_n)$ be two complete probability distributions then the information of kind β derived from the information function of kind β is in general given by (4) and if the information function f(x, y) of kind β is obtained in theorem I, then the corresponding information of kind β is

$$H_{\mathbf{n}}^{\beta}(P;Q) = \left[1 - \sum_{i=1}^{n} p_{i}\beta \ q_{i}^{1-\beta}\right] (1-\beta),^{-1} \beta \neq 1, \beta > 0.$$
(15)

Proof: Substituting the expression for f(x, y) from (8) in (4) we have

$$\begin{split} H^{\beta}_{n}(P;Q) &= (1-\beta)^{-1} \sum_{i=2}^{n} \left[P_{i}\beta \ Q_{i}^{1-\beta} - P_{i-1}^{\beta} \ Q_{i-1}^{1-\beta} - p_{i}\beta \ q_{i}^{1-\beta} \right] \\ &= (1-\beta)^{-1} \left[P_{n}\beta \ Q_{n}^{1-\beta} - P_{1}\beta \ Q_{1}^{1-\beta} - \sum_{i=2}^{n} p_{i}\beta \ q_{i}^{1-\beta} \right] \\ &= (1-\beta)^{-1} \left[1 - \sum_{i=1}^{n} p_{i}\beta \ q_{i}^{1-\beta} \right], \end{split}$$

clearly, it is (15).

PROPERTIES OF INFORMATION OF KIND &

Several properties like non-negativity, symmetry, null-information and expansibility⁹ can be easily derived for $H_n(P; Q)$.

We mention below some other properties for $H\beta_n$ (P; Q). Strongly Additive property:

where $P_{i}' = (p_{1i}, p_{2i}, \gg, p_{mi}), Q_{i}' = (q_{1i}, q_{2i}, \ldots, q_{mi}),$

$$p_1 * P_1' = (p_1 p_{11}, p_1 p_{21}, \ge, p_1 p_{m1})$$
 etc. and $\sum_{j=1}^m p_{ji} = 1, \sum_{j=1}^m q_{ji} = 1$

for all i = 1, 2, ... n.

An interesting special case of (16) is given below :

$$\overset{\beta}{H_{mn}} \left(\begin{array}{c} P^* P' \\ Q^* Q' \end{array} \right) = H^{\beta}_{n} \left(\begin{array}{c} P \\ Q \end{array} \right) + H^{\beta}_{m} \left(\begin{array}{c} P' \\ Q' \end{array} \right) + (\beta - 1) H^{\beta}_{n} \left(\begin{array}{c} P \\ Q \end{array} \right) H^{\beta}_{m} \left(\begin{array}{c} P' \\ Q' \end{array} \right)$$
(17)

where $P = (p_1, p_2, \ldots, p_n), \quad Q = (q_1, q_2, \ldots, q_n), \quad P' = (P_1, P_2, \ldots, P_m),$

$$Q' = (Q_1, Q_2, ..., Q_m)$$
 and $\sum_{j=1}^m P_j = 1$, $\sum_{j=1}^m Q_j = 1$.

When $\beta \rightarrow 1$, the last term in (17) vanishes and we have the well known additivity property of Kullback's information.

When $\beta \neq 1$ we have

$$H_{mn}^{\beta}\left(\begin{array}{c}P^{*}P'\\Q^{*}Q'\end{array}\right) < H_{\beta_{n}}\left(\begin{array}{c}P\\Q\end{array}\right) + H_{\beta_{m}}\left(\begin{array}{c}P'\\Q'\end{array}\right) \text{ for } \beta < 1.$$

ii Recursive-property

$$\begin{array}{l} \stackrel{\beta}{H_n} \left(\begin{array}{c} p_1, p_2, \dots, p_n \\ q_1, q_2, \dots, q_n \end{array} \right) - H_{\beta_{n-1}} \left(\begin{array}{c} p_1 + p_2, p_3, \dots, p_n \\ q_1 + q_2, q_3, \dots, q_n \end{array} \right) \\ \\ = (p_1 + p_2)^{\beta} (q_1 + q_2)^{1-\beta} H_2^{\beta} \left(\begin{array}{c} \frac{p_1}{p_1 + p_2}, \frac{p_2}{p_1 + p_2} \\ \frac{q_1}{q_1 + q_2}, \frac{q_3}{q_1 + q_2} \end{array} \right) \end{array}$$

with $p_1 + p_2$, $q_1 + q_2 > 0$. iii

$$H_m^{\beta}\left(\sum_{i=1}^n p_i q_{1i}, \dots, \sum_{i=1}^n p_i q_{mi}\right) \leq \sum_{i=1}^n p_i H_m^{\beta} \left(\frac{q_{1i}, \dots, q_{mi}}{q_1, \dots, q_m}\right) \text{ where } \sum_{j=1}^m q_{ji} = 1$$

$$Proof \quad H_{m}^{\beta}\left(\sum_{i=1}^{n} p_{i} q_{1i}, \dots, \sum_{i=1}^{n} p_{i} q_{mi}\right) = (1-\beta)^{-1} \left[1 - \sum_{j=1}^{m} \left(\sum_{i=1}^{n} p_{i} q_{ji}\right)^{\beta} q_{j}^{1-\beta}\right]$$

$$\leq (1-\beta)^{-1} \left[\sum_{i=1}^{n} p_i \left(1-\sum_{j=1}^{m} q_{ji} q_{j}^{1-\beta}\right) \right]$$
$$= \sum_{i=1}^{n} p_i H_m \beta \left(\frac{q_{1i}, \ldots, q_{mi}}{q_{1i}, \ldots, q_{mi}} \right).$$

because10

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$$\left(\sum_{i=1}^{n} p_{i} q_{ji}^{\beta}\right)^{\beta} > \sum_{i=1}^{n} p_{i} q_{ji}^{\beta} \text{ if } \beta < 1 (10)$$

CHARACTERIZATION OF INFORMATION IMPROVEMENT OF KIND \$

In this section we shall study a generalization of (2). Let the information improvement of kind β of the probability distributions P, Q and R be denoted by U_n be taken to satisfy the following postulates

DEF. SCI. J., VOL. 25, APRIL 1975

Postulate IV

$$U_{n}\left(\begin{array}{c}p_{1}, p_{2}, \dots, p_{n}\\q_{1}, q_{2}, \dots, q_{n}\\r_{1}, r_{2}, \dots, r_{n}\end{array}\right) = U_{n-1}\left(\begin{array}{c}P_{2}, p_{3}, \dots, p_{n}\\Q_{2}, q_{3}, \dots, q_{n}\\R_{2}, r_{3}, \dots, r_{n}\end{array}\right) + P_{2} Q_{2}^{1-\beta} R_{2}^{\beta-1} U_{2}\left(\begin{array}{c}p_{1}/P_{2}, p_{2}/P_{2}\\q_{1}/Q_{2}, q_{2}/Q_{2}\\r_{1}/R_{2}, r_{2}/R_{2}\end{array}\right)$$

where $P_3 = p_1 + p_2$, $Q_2 = q_1 + q_2$, $R_2 = r_1 + r_2$, P_2 , Q_2 , $R_2 > 0$.

Postulate V

 $U_{3}\left(\begin{array}{c}p_{1}, p_{2}, p_{3}\\q_{1}, q_{2}, q_{3}\\r_{1}, r_{2}, r_{3}\end{array}\right)$ is a symmetric function such that for any permutation of p's there are

the same permutations of q's and r's.

Postulate VI

$$U_{g}\left(\begin{array}{cc} \beta, \ 1-\beta\\ 1, \ 0\\ 1, \ 0\end{array}\right) = 1. \ (0 < \beta < 1).$$

Note IV: The postulate IV is a recursivity of information improvement of kind β it and it is a relation corresponding to additivity. Also postulate VI can be considered as a normalization property and here it includes the parameter β for the quantities where $0 < \beta < 1$.

Note V : If we take

$$f(x, y, z) = U_{2} \begin{pmatrix} x, 1-x \\ y, 1-y \\ z, 1-z \end{pmatrix},$$
(18)

then with the help of postulates IV and V, U_n (n > 3) can be expressed in terms of the single function f(z, y, z).

The postulate IV, at once gives the elegant form

$$U_n = \sum_{i=2}^n P_i Q_i^{1-\beta} R_i^{\beta-1} f(p_i / P_i, q_i / Q_i, r_i / R_i), \qquad (19)$$

where $P_i = p_1 + p_2 + \ldots + p_i$, $Q_i = q_1 + q_2 + \ldots + q_i$, $R_i = r_1 + r_2 + \ldots + r_i$; $i = 1, 2, \ldots, n$ with $P_n = Q_n = R_n = 1$.

When recursivity and symmetry of information improvement of kind β are applied to U_3 , we obtain the functional equation

$$f(x, y, z) + (1-x)(1-y)^{1-\beta}(1-z)^{\beta-1} f(u/1-x, v/1-y, w/1-z) = f(u, v, w) + (1-u)(1-v)^{1-\beta}(1-w)^{\beta-1} f(x/1-u, y/1-v, z/1-w)$$
(20)

for $x, y, z, u, v, w \in [0, 1]$ with $x + u, y + v, z + w \in [0, 1]$.

If we put z = x and w = u and take f(x, y, z) = f(x, y), (20) reduces to (5), the functional equation for information function of kind β and again if we put z = y = x and w = v = u and take f(x, x, x) = f(x), (20) reduces to Kendall's⁸ functional equation for information function.

We now adopt the following definition,

50

Definition: A real-valued solution f(x, y, z) of (20) defined on $[0, 1] \times [0, 1] \times [0, 1]$ is said to be an information improvement function of kind $\beta \in (0, 1)$, if it satisfies the following boundary conditions

$$f(0,0,0) = f(1,1,1),$$
(21)

$$f(\beta, 1, 1) = f(1 - \beta, 0, 0) = 1,$$
 (22)

for a given β such that $0 \leq \beta < 1$.

Again if f(x, y, z) is an information improvement function of kind $\beta \ (\neq 1)$ then the information improvement of kind $\beta \ (\neq 1)$ is given by (19).

Note VI: It is shown below that f(x,y,z) = f(1-x, 1-z) and that f(o,o,o) = o. Now (22) follows by putting $x = \beta$, y = 1, z = 1 in (20).

We now give characterizations of information improvement function and information improvement of kind β (\neq 1) in theorems III and IV respectively.

THEOREM III The only solution f(x, y, z) of (20) satisfying the additional conditions (21) and (22) is given by

$$f(x, y, z) = \begin{bmatrix} 1 - xy & -z \\ 1 - xy & -z \end{bmatrix} (1 - x) (1 - y)^{1 - \beta} (1 - z)^{\beta - 1} \int (1 - \beta)^{-1} (23)$$

for $(x, y, z) \in [0, 1] \times [0, 1] \times [0, 1] \sim (x, 1, 0) U(x, 0, 0) U(x, 0, 1) U(x, 1)$ (We shall take

Proof: Taking x = y = z = 0 in (20), (f 0, 0, 0) = 0 and hence from (21) we get f(1, 1, 1) = f(0, 0, 0) = 0.

Now replacing u, v, w, in (20) by 1 - x, 1 - y, and 1 - z respectively and applying (24) we have

$$f(x, y, z) = f(1 - x, 1 - y, 1 - z)$$
for x, y, z \epsilon [0, 1].
(25)

(24)

51

Let $p_1, p_2, q_1, q_2, r_1, r_2$, be six arbitrary numbers from the open interval (0, 1). Setting $p_1, = 1 - x$, $p_2 = u(1-x)^{-1}, q_1 = 1 - y, q_2 = v(1-y)^{-1}, r_1 = 1 - z, r_2 = w(1-z)^{-1}$ in (20) and using (25), we have

$$f(p_1, q_1, r_1) + p_1 q_1^{1-\beta} r_1^{\beta-1} f(p_2, q_2, r_2) + (1-p_1 p_2) (1-q_1 q_2)^{1-\beta} (1-r_1 r_2)^{\beta-1} \\ \cdot f(1-p_1/(1-p_1 p_2), 1-q_1/(1-q_1 q_2), 1-r_1/(1-r_1 r_2))$$
(26)

for $p_1, q_1, r_1 \in (0, 1)$, $p_2, q_2, r_2 \in [0, 1]$ such that $p_1 p_2 \neq 1$, $q_1 q_2 \neq 1$, $r_1 r_2 \neq 1$. Consider the function

$$F(p_1, p_2; q_1, q_2; r_1, r_2) = f(p_1, q_1, r_1) + \begin{bmatrix} 1 - \beta & \beta - 1 \\ p_1 q_1 & r_1 \end{bmatrix} + (1 - p_1) (1 - q_1)^{1 - \beta} (1 - r_1)^{\beta - 1} \end{bmatrix} \begin{pmatrix} \beta - 1 \\ r_1 \end{pmatrix} f(p_2, q_2, r_2), p_1, p_2, q_1, q_2, r_1, r_2 \in (0, 1).$$
(27)

We shall show that $F(p_1, p_2; q_1, q_2; r_1, r_2)$ is symmetric, i.e.,

$$F(p_1, p_2; q_1, q_2; r_1, r_2) = F(p_2, p_1; q_2, q_1; r_2, r_1).$$
(28)

Let us take $\beta \neq 1$ ($\beta > 0$), then by (26),

$$F \ p_1, \ p_2; \ q_1, \ q_2; \ r_1, \ r_2) = f \ (p_1 \ p_2, \ q_1 \ q_2, \ r_1 \ r_2) + (1 - p_1 \ p_2) \ (1 - q_1 \ q_2)^{1 - \beta} (1 - r_1 \ r_2^{\beta - 1}) \\ \left[f \left(1 - p_1 | (1 \ p_1 \ p_2), \ 1 - q_1 | (1 - q_1 \ q_2), \ 1 - r_1 | (1 - r_1 \ r_2) \right) + (1 - p_1 | (1 - p_1 \ p_2)) \ (1 - q_1 | (1 - q_1 \ q_2)^{1 - \beta} \right] \right]$$

$$(1 - r_1/(1 - r_1 r_2)) f (p_2, q_2, r_2)$$
(29)

Again setting
$$p^{r_1} \equiv (1 - p_1) (1 - p_1 p_2)^{-r_1} q^{r_1} \equiv (1 - q_1) (1 - q_1 q_2)^{-r_1} r_1 \equiv (1 - r_1) (1 - r_1 r_2)^{-r_1}$$
,

$$A (p_1, p_2; q_1, q_2, r_1, r_2) = = f (1 - p_1) (1 - p_1 p_2) \cdot 1 - q_1 | (1 - q_1 q_2), 1 - r_1 | (r - r_1 r_2)) + (1 - p_1 | 1 - p_1 p_2) (1 - q_1 | 1 - q_1 q_2)^{1 - \beta} (1 - r_1 | 1 - r_1 r_2).^{\beta - 1} f (p_2 q_2 r_2)$$

$$= f (p^{r_1}, q^{r_1}, r^{r_1}) + p^{r_1} q^{r_1} f^{r_1} f (p_2, q_2, r_2)$$

$$= f (1 - p_2 | (1 - p_1 p_2), 1 - q_2 | (1 - q_1 q_2), 1 - r_2 | (1 - r_1 r_2)) + (1 - p_2 | 1 - p_1 p_2) (1 - q_2 | 1 - q_1 q_2)^{1 - \beta} (1 - r_2 | 1 - r_1 r_2)^{\beta - 1} f (p_1, q_1, r_1)$$

$$= A (p_2, p_1; q_2, q_1; r_2, r_1)$$

Then it follows from (27) that

$$F(p_1, p_2; r_1, q_2;) q_1, r_2) - F(p_2, p_1; q_2, q_1; r_2, r_1) = 0$$
, proving (28)

Next putting $p_2 = \beta$, $q_2 = 1$, $r_2 = 1$ or $p_2 = 1 - \beta$, $q_2 = 0$, $r_2 = 0$ in (28) and using the definition of $F(p_1, p_2; q_1, q_2; r_1, r_2)$ and (22) we get

$$0 = f(p_1, q_1, r_1) + \left[p_1 q_1 \stackrel{1-\beta}{r_1} r_1 \stackrel{\beta-1}{r_1} + (1-p_1) (1-q_1) \stackrel{1-\beta}{r_1} (1-r_1) \stackrel{\beta-1}{r_1} \right] - 1 - \beta f(p_1, q_1, r_1)$$

From this it follows that

$$f(p_1, q_1, r_1) = \left[\begin{array}{c} 1 - p_1 q_1 & \beta - 1 \\ r_1 & -(1 - q_1) (1 - q_1) & (1 - r_1) \end{array} \right] (1 - \beta)^{-1}$$

for all $p_1, q_1, r_1 \in (0, 1).$

The result is true even when $p_1 = q_1 = r_1 = 0$ or $p_1 = q_1 = r_1 = 1$.

THEOREM IV Let $P = (p_1, p_2, \ldots, p_n), Q = (q_1, q_2, \ldots, q_n)$ and $R = (r_1, r_2, \ldots, r_n)$ be three complete probability distributions then the information of kind β derived from the information improvement function of kind β is in general given by (19) and if the information improvement function f(x, y, z)of kind β is as obtained in theorem III then the corresponding information improvement of kind β is

$$H_{n}^{\beta}(P; Q; R) = \left[1 - \sum_{i=1}^{n} p q_{i}^{1-\beta} r_{i}^{\beta-1}\right] (1-\beta)^{-1}, \ \beta \neq \beta > 0.$$
(30)

proof: Substituting the expression for f(x, y, z) from (23) in (19) we have

$$\begin{split} H_{n}^{\beta}(P;Q;R) &= (1-\beta)^{-1} \sum_{i=2}^{n} \left[P_{i} Q_{i}^{1-\beta} R_{i}^{\beta-1} P_{i-1} Q_{i-1}^{1-\beta} R_{i-1}^{\beta-1} - p_{i} q_{i}^{1-\beta} r_{i}^{\beta-1} \right] \\ &= (1-\beta)^{-1} \left[P_{n} Q_{n}^{1-\beta} R_{n}^{\beta-1} - P_{1} Q_{1}^{1-\beta} R_{1}^{\beta-1} - \sum_{i=2}^{n} p_{i} q_{i}^{1-\beta} r_{i}^{\beta-1} \right] \\ &= (1-\beta)^{-1} \left[1 - \sum_{i=1}^{n} p_{i} q_{i}^{1-\beta} r_{i}^{\beta-1} \right], \\ &\quad \text{clearly, which is (30).} \end{split}$$

PROPERTIES OF INFORMATION IMPROVEMENT OF KIND &

Several properties like symmetry, null-information, expansibility can be easily derived for H_n (P;Q;R) β We mention below, some other properties for H_n (P;Q;R).

52

SONI : Functional Equation in Information Theory

(i) Strongly (Non-Commutative Additive Property).

$$H_{mn}^{\beta} = \begin{pmatrix} p_{1} P'_{1}, p_{2} P'_{2}, \dots, p_{n} P'_{n} \\ q_{1} Q'_{1}, q_{2} Q'_{2}, \dots, q_{n} Q'_{n} \\ r_{1} R'_{1}, r_{2} R'_{2}, \dots, r_{n} R'_{n} \end{pmatrix} = \\ = H_{n}^{\beta} \begin{pmatrix} p_{1}, p_{2}, \dots, p_{n} \\ q_{1}, q_{2}, \dots, q_{n} \\ r_{1}, r_{2}, \dots, r_{n} \end{pmatrix} + \sum_{i=1}^{n} p_{i} q_{i} r_{i} H_{m}^{\beta} \begin{pmatrix} P'_{i} \\ Q'_{i} \\ R'_{i} \end{pmatrix}$$
(31)
$$P'_{i} = (p_{1i} p_{2i}, \dots, p_{n}) \cdot Q'_{i} = (q_{i}, q_{i}, \dots, q_{m}),$$

where

$$\begin{array}{c} r_{i} = (p_{1i}, p_{2i}, \ldots, p_{mi}), \ q_{i} = (q_{1i}, q_{2i}, \ldots, q_{mi}), \\ R'_{i} = (r_{1i}, \ldots, r_{mi}), \ p_{1}^{*} P'_{1} = (p_{1} \ p_{11}, p_{1} \ p_{21}, \ldots, p_{1} \ p_{mi}) \quad \text{etc.}, \\ \\ \frac{\Sigma}{j=1} p_{ji}, \ldots, = 1, \ \frac{\Sigma}{j=1} q_{ji} = 1 \ \frac{\Sigma}{j=1} r_{ji} = 1 \ \text{for all} \ i = 1, 2, \ldots, n. \end{array}$$

and

An interesting special case of (31) is given below.

$$\overset{\beta}{\underset{\substack{H_{mn}}}{}} \begin{pmatrix} P^* P' \\ Q^* Q' \\ R^* P' \end{pmatrix} = \overset{\beta}{\underset{\substack{H_n}}{}} \begin{pmatrix} P \\ Q \\ R \end{pmatrix} + \overset{\beta}{\underset{\substack{H_m}}{}} \begin{pmatrix} P' \\ Q' \\ R' \end{pmatrix} + (\beta - 1) \overset{\beta}{\underset{\substack{H_n}}{}} \begin{pmatrix} P \\ Q \\ R \end{pmatrix} \overset{\beta}{\underset{\substack{H_m}}{}} \begin{pmatrix} P' \\ Q' \\ R' \end{pmatrix} (32)$$

$$P = (p_1, p_2, \dots, p_n) \text{ and } P' = (P_1, P_2, \dots, P_m) \text{ etc. and }$$

where

$$\sum_{j=1}^{m} P_j = 1, \sum_{j=1}^{n} Q_j = 1 \text{ and } \sum_{j=1}^{m} R_j = 1.$$

When $\beta \rightarrow 1$, the last term in (32) vanishes and we have the known additive property of Theil's information-improvement.

When $\beta \neq 1$, we have

$$\begin{array}{c} \beta\\ H_{mn} \\ \begin{pmatrix} P^* P\\ \mathring{Q} Q'\\ R^* R' \end{pmatrix} \\ & \geq H_n \\ \begin{pmatrix} P\\ Q\\ R \end{pmatrix} \\ & + H_m \\ \begin{pmatrix} P'\\ Q'\\ R' \end{pmatrix} \\ & \begin{pmatrix} P'\\ Q'\\ R' \end{pmatrix} \\ & \begin{pmatrix} \beta -1 \end{pmatrix} \\ & \begin{pmatrix} P\\ Q\\ R \end{pmatrix} \\ & \begin{pmatrix} P\\ Q\\ R \end{pmatrix} \\ & \begin{pmatrix} P'\\ H_m \\ R' \end{pmatrix} \\ & \begin{pmatrix} P'\\ Q'\\ R' \end{pmatrix} \\ & \geq 0. \end{array}$$

according as

(ii) Recursive-Property

2

with $p_1 + p_2, q_1, + q_2, r_1 + r_2 > 0$.

$$\overset{\beta}{\underset{i=1}{\overset{\mu}{\underset{i=1}{\overset{m}{\atop}}}}} \left(\begin{array}{c} \overset{\Sigma}{\underset{i=1}{\overset{p_{i}}{\atop}}} q_{1i}, \ldots, \overset{\Sigma}{\underset{i=1}{\overset{p_{i}}{\atop}}} q_{i} q_{mi} \\ q_{1}, \ldots, q_{m} \\ q_{1}, \ldots, q_{m} \end{array} \right) = \begin{array}{c} \overset{n}{\underset{i=1}{\overset{\beta}{\atop}}} p_{i} H_{m} \\ \overset{\beta}{\underset{i=1}{\overset{m}{\atop}}} \left(\begin{array}{c} q_{1i}, \ldots, q_{mi} \\ q_{1}, \ldots, q_{mi} \\ q_{1}, \ldots, q_{mi} \end{array} \right)$$

ACKNOWLEDGEMENTS

The author wishes to express his sincere thanks to Dr. Ahu Sharma, Reader in Mathematics, University of Delhi for guidance and encouragement in carrying out this research work and to Professor U. N. Singh, Dean, Faculty of Mathematics for providing facilities in the department.

Thanks are also due to C.S.I.R. for the award of Research Fellowship.

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